

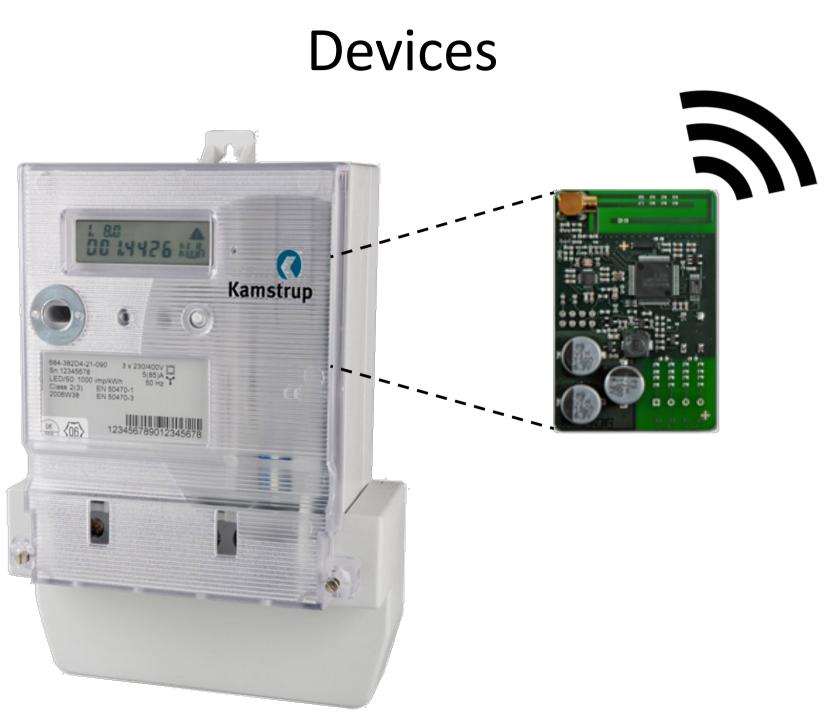
Distributed Computing and Systems



# ICT Support for Adaptiveness and (Cyber)security in the Smart Grid

#### Low-Power Wireless Communication

Olaf Landsiedel



### Routing



#### Lifetime



# **Goal for Today**

- Devices
  - Low-power wireless communication
  - Or a "wireless sensor node" in general
- Routing in low-power wireless networks

   From source to sink
- Energy efficiency: often battery driven
  - Enable a life-time of years: allow devices to sleep

Low-Power Wireless

#### **Devices**

## Requirements for such a Device

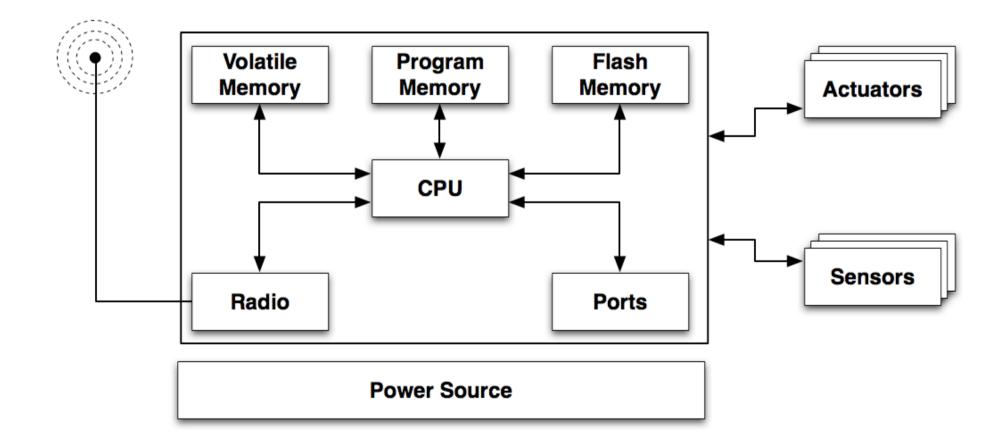
Low cost

• High energy efficiency

• Small size



#### A Sensor Node (or low-power wireless device)



#### A Sensor Node (or low-power wireless device)

- TI MSP 430 (16 bit RISC)
  - 8 MHz
  - 10 KB RAM
  - 48 KB code, 1MB flash
- Chipcon CC2420 radio
  - IEEE 802.15.4 compliant
  - 50 m. range indoor,250 m. range outdoor
  - bandwidth 250 kbits/s
- On-board antenna



# Why not use WiFi or Bluetooth?

- WiFi/WLAN (IEEE 802.11)
  - Topology: mesh + single hop
  - Throughput: >100 Mbps
  - Power Consumption: ~300mW
- Bluetooth
  - Topology: Single-hop network
    - Master <-> Slave
    - Not good for multi-hop networking
  - Throughput: up to 24 Mbit/s
  - Power Consumption: up to 30mW





#### Summary: Device

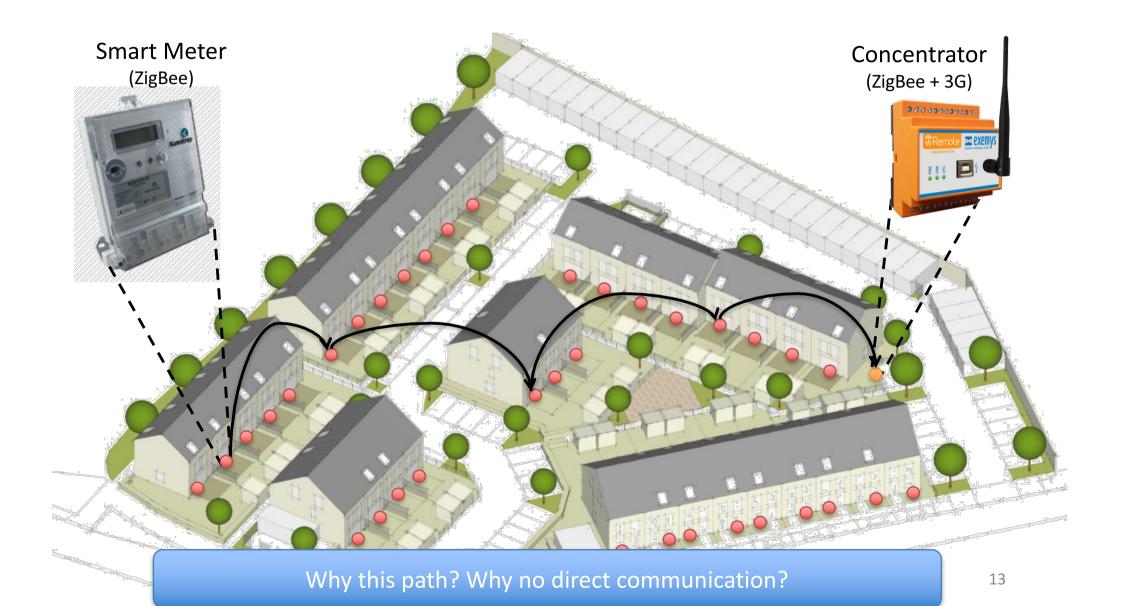
- Low-Power Hardware
  - Simple Processor
  - Simple, energy-efficient radio

• Low cost, low energy consumption

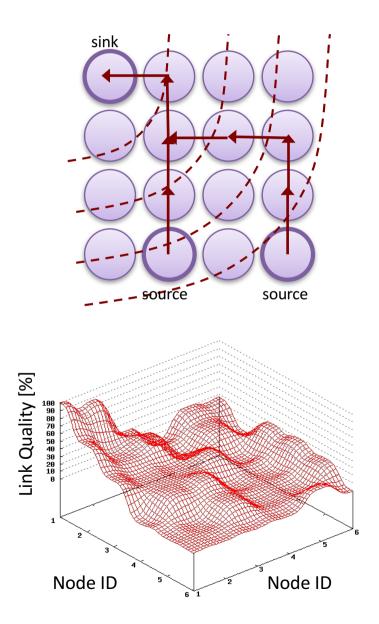
Low-Power Wireless

#### Routing

### Routing



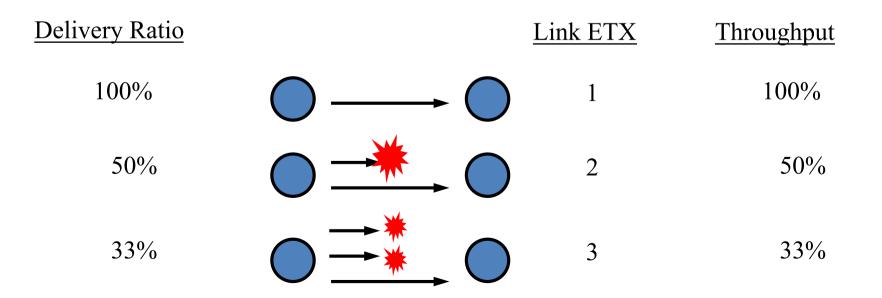
# **Routing Metrics**



- Path Selection
  - Which path to select?
  - Routing Metric?
    - Minimize Hops?
    - Reliability?
- Wireless Links — Highly dynamic

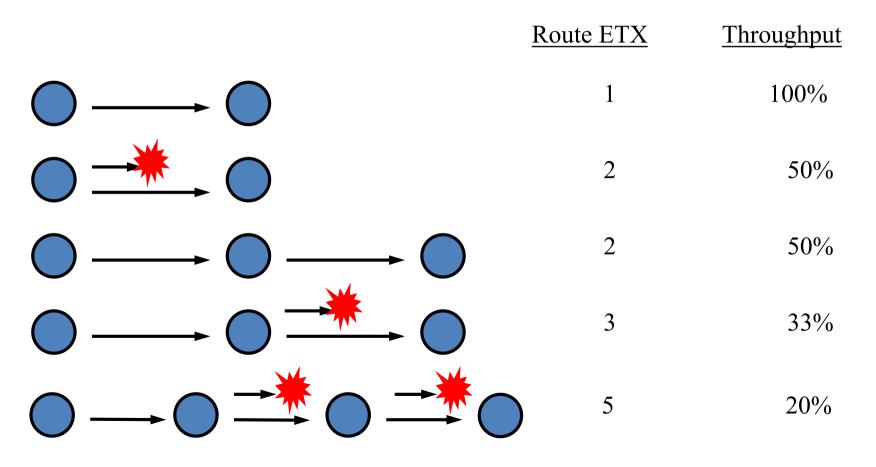
## Routing Metric: ETX

- Goal: Minimize total transmissions per packet
  - Use Metric: Expected Transmission Count (ETX)
    - Measure link over a time to determine ETX
  - Link throughput  $\approx 1$  / Link ETX

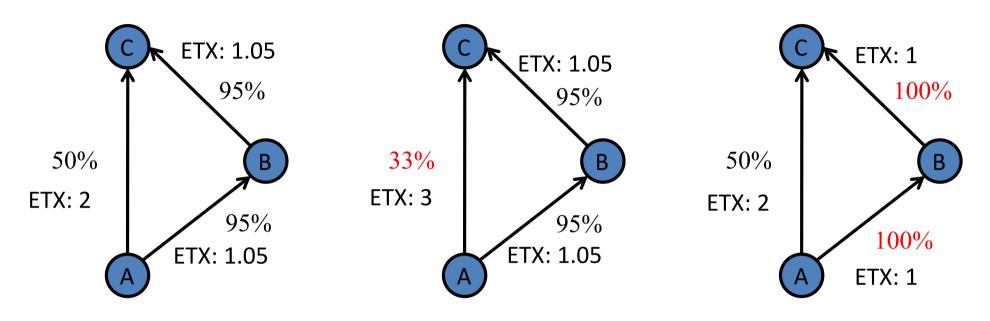


#### Route ETX

- Route ETX = Sum of link ETXs
  - Communication is expensive
    - ETX predicts the tx count of a packet -> Reflects energy
  - Route selection:
    - Choose route with lowest route ETX



## Question: Which Route is better?



- Which route to take
  - − A -> B −> C or A -> B?
  - Example 1
    - A -> B -> C: 2.1 TX
    - A -> B: 2 TX
    - Take A->B

- Example 2
  - A -> B -> C: 2.1 TX
  - A -> B: 3 TX
  - Take A -> B -> C
- Example 3
  - A -> B -> C: 2 TX
  - A -> B: 2 TX
  - Take any

#### Summary: Routing Metric



- Expected Transmission Count (ETX)
  - Minimize total number of transmission
  - Good for energy: More transmissions -> more energy
  - Combines hops and reliability into single metric

Low-Power Wireless

#### **Sleeping Devices**

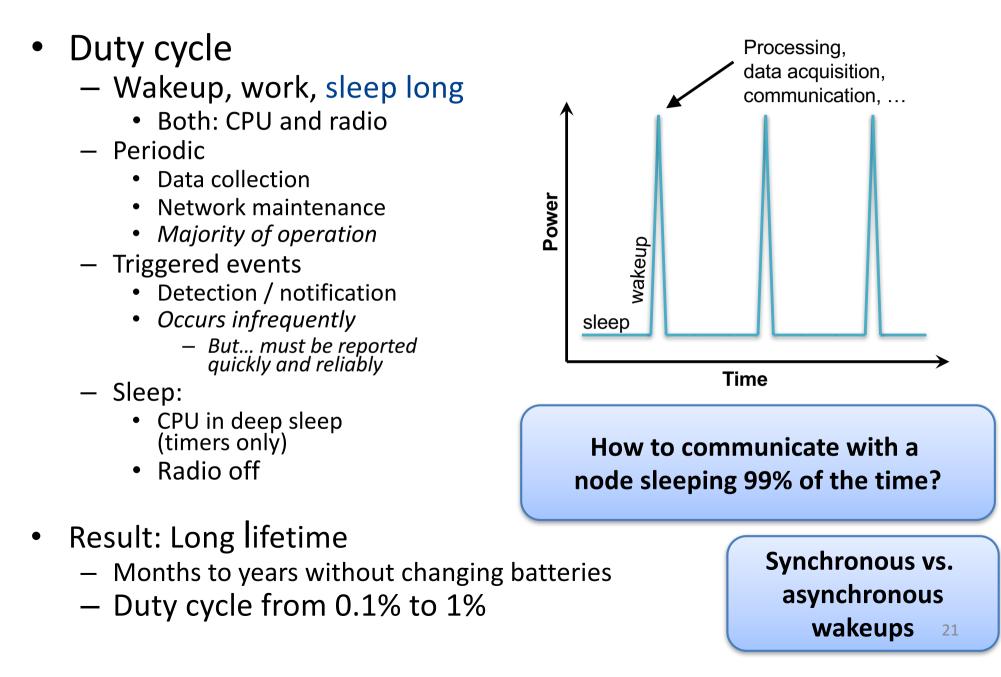
# Energy-Efficient MAC

- Targeted life time of WSN
  - Months or years
- Simple back of the envelope calculation:
  - AA battery: About 2000 mAh
  - CC2420 radio:
     19.7mA in RX mode
     (listening to channel)
  - 2000mAh / 19.7mA = 101.5 hours
    - = 101.5 hour
    - = 6 days

Operation	Telos
Minimum Voltage	1.8V
Mote Standby (RTC on)	5.1 µA
MCU Idle (DCO on)	54.5 $\mu A$
MCU Active	1.8 mA
MCU + Radio RX	21.8 mA
MCU + Radio TX (0dBm)	19.5 mA
MCU + Flash Read	4.1 mA
MCU + Flash Write	15.1 mA
MCU Wakeup	6 µs
Radio Wakeup	$580~\mu{ m s}$

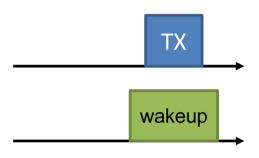
We want month or years: How?
 →Keep radio off most of the time

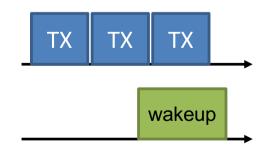
# Solution: Duty Cycling



# **Duty Cycling**

- Synchronous duty cycling
  - Knowing the wakeup time of destination
    - Transmit accordingly
  - Advantage: very energy efficient
  - Disadvantage: requires synchronization
- Asynchronous duty cycling
  - Not knowing the wakeup time
    - Example: Repeat transmission until destination wakes up and acknowledges
  - Advantage: simple, no time synchronization
  - Disadvantage: not as energy efficient



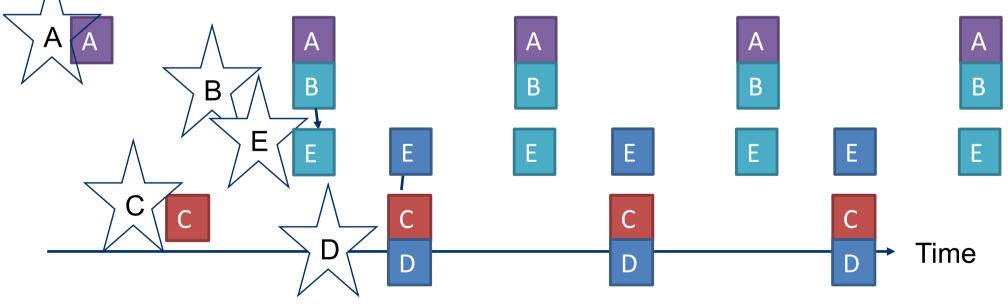


# Synchronous Duty Cycling

- Idea:
  - Switch nodes, radios off
  - Ensure that neighboring nodes turn on simultaneously
    - To allow packet exchange (rendezvous)
    - Requires Time Synchronization
    - Called "Synchronous duty cycling"
- In wakeup phase
  - Only in these *active periods*, packet exchanges happen
  - Need to also exchange wakeup schedule between neighbors

# Synchronous Duty Cycling

- Nodes try to pick up schedule synchronization from neighboring nodes
- If no neighbor found, nodes pick some schedule to start with
- If additional nodes join, some node might learn about two different schedules from different nodes
  - "Synchronized islands"
- To bridge this gap, it has to follow both schemes

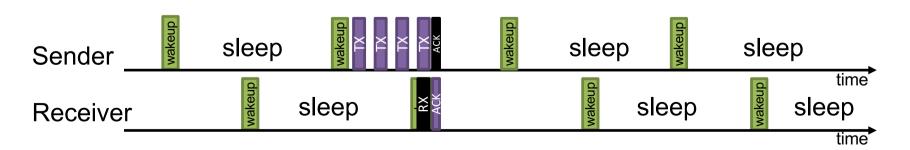


#### Synchronous Duty Cycling: Discussion

- Pro: Energy-Efficient
  - A node sleeps most of the time
  - Periodically wake up for short intervals to see if any node is transmitting a packet
- Cons
  - Time sync overhead
    - Account for clock drifts etc.
    - Add guard spaces
  - Some nodes are in multiple "clusters"
    - More wakeups
    - Have higher energy consumption

# Asynchronous Duty Cycling

- So far: Periodic sleeping
  - Need some means to synchronize wake up of nodes
    - Ensure rendezvous between sender and receiver
- Alternative option: Don't try to explicitly synchronize nodes
  - Have receiver sleep and only periodically sample the channel
- Repeat packet until receivers wakes up
  - And acknowledges
  - No Synchronization required! Asynchronous duty cycle



#### Asynchronous Duty Cycling: Discussion

- Pro: Energy-Efficient
  - A node sleeps most of the time
  - No need for time sync
  - Periodically wake up for short intervals to see if any node is transmitting a packet
- Cons
  - Transmission are costly
    - Especially when nodes wakeup rarely
  - A single transmissions is repeated many times
    - High channel utilization in this time

## Summary

- Devices: cheap, low-power
  - Low-power wireless
- Routing: Expected Transmission Count (ETX)
   Account for link dynamics
- Synchronous and asynchronous duty cycling
  - Sleeping devices

#### Questions?

In part, inspired from / based on slides and figures from Jochen Schiller, Holger Karl, Klaus Wehrle, Kyoung-Don Kang, Leonardo Leiria Fernandes, Joe Polastre, Chenyang Lu, Leo Selavo, Luca Mottola, Adam Dunkels, and many others